

# Optimal Placement of Retry-Based Fault Recovery Annotations in HPC Applications



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#### Motivation

- Mean-time-between-failures may increase in future HPC systems
- We may not rely only on the checkpoint/restart approach
- Retry-based methods allow recovery by re-trying a code region
- Two retry-based approaches:

(1) Identify or create idempotent code (i.e., code where reexecution is free of side effects)

(2) Place retry annotations (involves micro-checkpointing data) However, how do we place optimally these annotations?

## **Overview of Existing Fault Recovery Methods**

Replication in Space

- Redundant multi-threading
- Redundant VMs
- Lockstepping

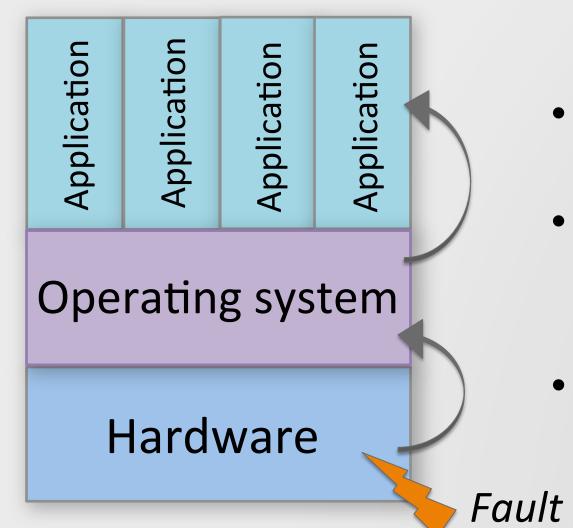
Incur hardware overheads
Inefficient for parallel codes

Replication in *Time* 

- Checkpoint / Restart
- Retry idempotent code
- Micro-checkpoint / Retry

## Assumptions for the Retry-Based Model

Replication in Time



- Fault detection in hardware
- Notifications of recoverable faults (e.g., via synchronous traps)
- Recovery at application level

## **Annotations are Required to Express Recovery Scope**

- Programmer annotates (or protect) code block
- If a fault occurs, a code block is re-executed
- The decision on where to place annotations is critical or large overheads can occur
- No research work has evaluated how to optimally place these annotations

#### Original code

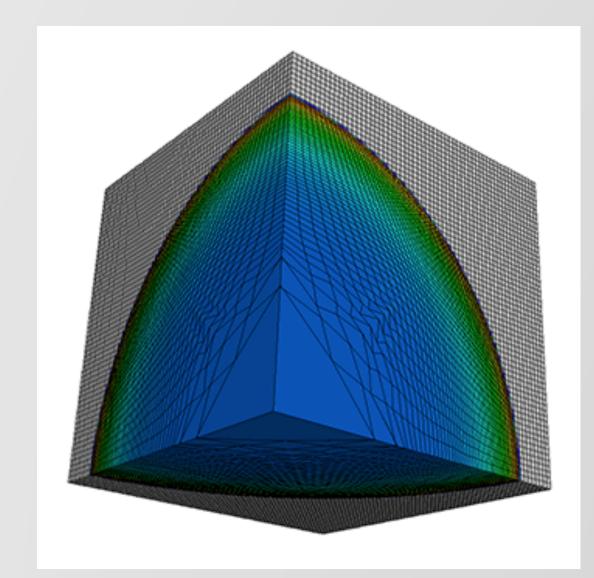
```
void function(double *array)
{
  for (...)
   array[i] = array[i-1] + ...
}
```

# Annotated code

```
void function(double *array)
{
   RETRY{
    for (...)
     array[i] = array[i-1] + ...
}
}
```

# Sample Application: LULESH

Shock hydrodynamics problem



How do we annotate the code?

main() {
 /\* init code...\*/
 while() {
 funct1();
 funct2();
 funct3();
 }
}

